

# QCI

## This class

1. Intro to QI
2. QECC
3. QIT
4. q algorithms
5. complexity
6. physics

## Why?

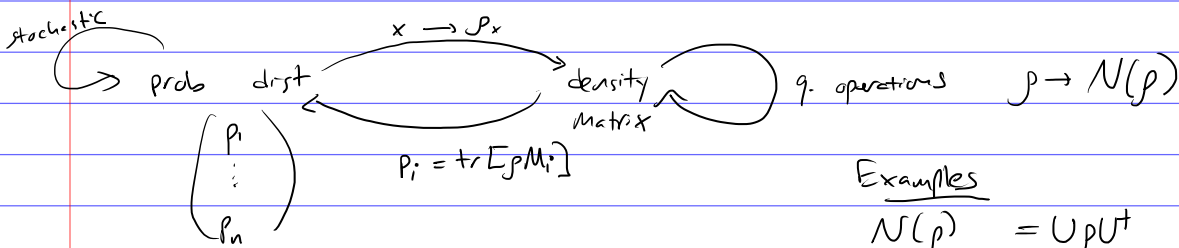
- Building and using a QC
- Using QI for tasks like sensing & control
- understanding QI in nature

## States and operations

States	deterministic	random
classical	$C^d$	$p \in \mathbb{R}^d, p \geq 0, \ p\ _1 = 1$
quantum	$ \psi\rangle \in C^d$ $\langle\psi \psi\rangle = 1$	$\rho \in C^{d \times d}, \rho \geq 0, \text{tr} \rho = 1$

operations	deterministic	random
	$f: C^d \rightarrow C^d$ $U(C^d)$	$T \in \mathbb{R}^{d \times d}$ stochastic q. operations / channels / TPCP maps

we can also map between distributions and states



$$M_1 + \dots + M_n = I$$

$$M_i \geq 0$$

## Examples

$$N(\rho) = U\rho U^\dagger \quad \text{depolarizing}$$

$$N(\rho_{AB}) = \text{tr}_B[\rho_{AB}] \quad \text{depolarizing}$$

$$N(\rho) = \rho \otimes \sigma \quad \text{amp. damping}$$

$$N(\rho) = V\rho V^\dagger \quad V \text{ isometry}$$

$$N(\rho) = \sum_i p_i N_i(\rho)$$

$$N(\rho) = \sigma$$

$$N(\rho) = \sum_i \text{tr}[\rho M_i] \sigma_i$$

i "measure-and-prepare"

# Three ways to characterize q. operators

- 1) physical (Stinespring)
- 2) algebraic (Kraus)
- 3) axiomatic (TPCP)

① "one equation to rule them all"

$$\frac{d}{dt} |\Psi_{\text{universe}}\rangle = \frac{-iH}{\hbar} |\Psi_{\text{universe}}\rangle$$

- everything is unitary
- partial trace is subjective

possible operations within this picture

$\rho \rightarrow \rho \otimes \sigma$  suffices to use  $\rho \rightarrow \rho \otimes |\alpha\rangle\langle\alpha|$

$\rho \rightarrow U\rho U^\dagger$

$U(\rho \otimes |\alpha\rangle\langle\alpha|)U^\dagger = V\rho V^\dagger$  for an isometry  $V = U(I \otimes |\alpha\rangle)$

$J_{SE} \rightarrow \text{tr}_E \rho$

$N(\rho) = \text{tr}_E V\rho V^\dagger$        $V: S \rightarrow B \otimes E$   
or  $V: A \rightarrow B \otimes E$

②

$V: A \rightarrow B \otimes E$

means  $V = \sum_e V_e \otimes |e\rangle$  where  $\{|e\rangle\}$  is an o.n. basis for  $E$   
 $V_e \in L(A, B) = \text{lin. operators from } A \rightarrow B$

$N(\rho) = \text{tr}_E \sum_{e_1, e_2} V_{e_1} \rho V_{e_2}^\dagger \otimes |e_1\rangle\langle e_2|$

$N(\rho) = \sum_e V_e \rho V_e^\dagger$

$\{V_e\}$  are Kraus operators      Kraus decomposition or OSR  
operator-sum representation

$V = \begin{pmatrix} V_1 \\ V_2 \\ \vdots \\ V_{d_E} \end{pmatrix}$

$V_i$  is  $d_B \times d_A$   
constraints?

$V$  isometry means  $V^\dagger V = I = \sum_{e_1} V_{e_1}^\dagger \otimes \langle e_1| \sum_{e_2} V_{e_2} \otimes |e_2\rangle = \sum_e V_e^\dagger V_e$

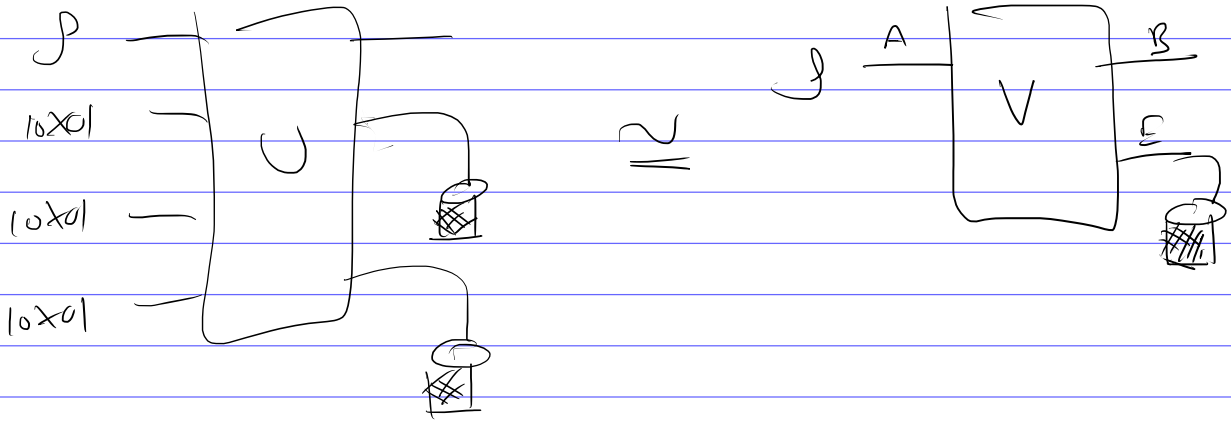
(\*)  $\sum_e V_e^\dagger V_e = I$       Kraus operator conditions.

conversely if  $\{V_e\}$  satisfies (\*) then  $V = \sum_e V_e \otimes |e\rangle$  is an isometry and  $N$  is a physical map

examples 1) unitary

2) measurement  $\{M_i\} \rightarrow V = \sum |i\rangle \otimes \sqrt{M_i}$   
 $\text{tr}_E V \rho V^\dagger = \sum_i \text{tr}(M_i \rho) |i\rangle \langle i|$

composition?



or OSR  $\rho \rightarrow \sum_e V_e \rho V_e^\dagger \rightarrow \sum_{e,f} W_{ef} V_e \rho V_e^\dagger W_{ef}^\dagger = (W \rho V)^\dagger$

TPCP maps

$V$  vector space

$\rho \in L(V) = L(V, V)$

$N \in L(L(V)) = L(L(V), L(V))$

i.e.  $N$  is linear, and Hermiticity preserving

1. TP  $\text{tr } N(\rho) = \text{tr } \rho$

2. CP  $\rho \geq 0 \Rightarrow N(\rho) \geq 0$

and  $(N \otimes \text{id})(\rho) \geq 0$

why necessary?  $\rho \geq 0 \Rightarrow \rho^T = T(\rho) \geq 0$  but  $(T \otimes \text{id})(\frac{F}{d}) = \frac{F}{d}$

$N(\rho) = \sum_e V_e \rho V_e^\dagger$   $\text{tr } N(\rho) = \text{tr } \rho \underbrace{\sum_e V_e^\dagger V_e}_I$

$\rho = W W^\dagger$   $N(\rho) = \sum_e V_e W W^\dagger V_e^\dagger \geq 0$

Different channel pictures are based on idea of purifications

$\forall \rho_A \exists$  purification  $\psi_{AB}$  s.t.  $\rho_A := \text{tr}_B[\psi] = \rho_A$   
 $= |\psi\rangle\langle\psi|$

recall

density matrices

$$D_d = \{ \rho \in \mathbb{C}^{d \times d} : \rho \geq 0, \text{tr} \rho = 1 \}$$

$$\rho \geq 0 \Leftrightarrow \rho = C^T C \Leftrightarrow \text{eigs}(\rho) \geq 0 \Leftrightarrow \langle v | \rho | v \rangle \geq 0 \quad \forall |v\rangle$$

purifications

$$|\psi\rangle_{AB} = \sum_{ij} C_{ij} |i\rangle |j\rangle = \text{vec}(C) \quad \psi_A = \sum_{ij} C_{ij} C_{ij}^* |i\rangle \langle i| = C C^T$$

given  $\rho_A$  can find  $|\psi\rangle_{AB}$  s.t.  $\psi_A = \rho_A$

What are the possible purifications?

can understand in terms of SVD

$$C = U D V^T$$

$$D = \text{diag}(\lambda_1, \lambda_2, \dots) \quad \lambda_1 \geq \lambda_2 \geq \dots \geq 0$$

$$\rho = C C^T = U D^2 U^T$$

eig( $\rho$ ) determines  $D$

$U$  is determined up to  $UR$  s.t.  $[R, D] = 0$

$V$  is arbitrary

$$|\psi\rangle = \sum_{ij} C_{ij} |i\rangle \otimes |j\rangle$$

$$(A \otimes B) |\psi\rangle = \sum_{ijk} A_{ki} C_{ij} B_{je}^T |k\rangle |e\rangle = \text{vec}(A C B^T)$$

Then given  $|\psi\rangle_{AB}$ ,  $|\chi\rangle_{AB}$  with  $\psi_A = \chi_A$

$$\exists \text{ unitary } W \text{ s.t. } (\mathbb{I} \otimes W) |\psi\rangle = |\chi\rangle$$

pf  $\Leftarrow$   $\mathbb{I} \otimes W$  doesn't change  $A$  marginal

$$\Rightarrow \begin{aligned} |\psi\rangle &= \text{vec}(C_1) & C_1 &= U_1 D_1 V_1^T & W &\text{ can remove } V_1^T, U_2^T \\ |\chi\rangle &= \text{vec}(C_2) & C_2 &= U_2 D_2 V_2^T \end{aligned}$$

$$C_1 C_1^T = C_2 C_2^T$$

$$U_1 D_1^2 U_1^T = U_2 D_2^2 U_2^T \Rightarrow D_1 = D_2 =: D$$

$$U_1 R = U_2 \text{ for some } R \text{ s.t. } [R, D] = 0$$

complete proof on pset

Cor  $|\psi\rangle_{AB}$ ,  $|\chi\rangle_{AB}$  have  $\psi_A = \chi_A$

$\Downarrow$

either  $\exists$  isometry  $W$  s.t.  $(\mathbb{I} \otimes W) |\psi\rangle = |\chi\rangle$  or  $(\mathbb{I} \otimes W) |\chi\rangle = |\psi\rangle$

# info theoretic crypto

OKO

coin-flipping: output random bit with  $\Pr[E] = p \approx 1/2$

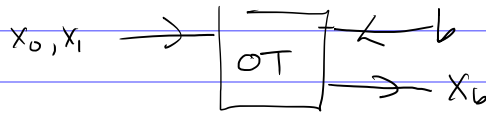
NO  
YES

strong  
weak

$$p \approx 1/2$$

Alice can choose  $p \leq \frac{1}{2} + \epsilon$   
Bob  $p \geq \frac{1}{2} - \epsilon$

oblivious transfer



## Bit commitment

Alice has input  $b$ .

commit phase

Bob can't guess  $b$  (hiding)

reveal phase

Bob outputs  $b$  or REJECT (binding)

$\Pr[\text{REJ}] \approx 0$  for honest players (valid)

OT  $\rightarrow$  BC  $\rightarrow$  strong coin flipping  
Thus info-theoretic secure q. BC is impossible.

can understand this through perfectness

## No B.C.

recall channels ① TPCP ② Kraus ③  $N(p) = \frac{1}{\epsilon} \log p^T$

- purify protocol

- cheaters can access discarded systems

$|\psi_b\rangle_{AB}$  after commit. hiding  $\Rightarrow \psi_0^B = \psi_1^B$   
 $\Rightarrow$  not binding

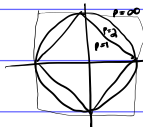
# Metrics on states and prob. distributions

## Norms $\|\cdot\|$

- satisfying
- 1)  $\|cv\| = |c| \cdot \|v\|$  for  $c \in \mathbb{C}$
  - 2)  $\|v+w\| \leq \|v\| + \|w\|$
  - 3)  $\|v\| = 0 \Leftrightarrow v=0$

$\|v\|^2 = \langle v, v \rangle$  only for 2-norm

$L_p$   $x \in \mathbb{C}^d$   $\|x\|_{L_p} = \|x\|_p = \left( \sum_{i=1}^d |x_i|^p \right)^{1/p}$



- $L_1$   $\sum |x_i|$
- $L_2$  Euclidean
- $L_\infty$   $\max_i |x_i|$

$S_p =$  Schatten-p  $\|X\|_{S_p} = \|X\|_p = \| \text{svals}(X) \|_{L_p}$

$\|X\|_{S_1} = \text{tr} \sqrt{X^\dagger X} = \text{tr} |X|$  if  $X$  is normal

$\|X\|_{S_\infty} = \text{op. norm} = \text{biggest sval} = \|X\|$

unit sphere/ball  $S(\cdot), B(\cdot)$

pure state QM  $S(L_2)$

probability dists  $S(L_1) \cap \mathbb{R}_+^d$

density matrices  $S(S) \cap \text{psd}$

measurement ops  $B(S_\infty)$

$L_p$  and  $L_q$  dual if  $\frac{1}{p} + \frac{1}{q} = 1$   
 $S_p$  ;  $S_q$  " " " "  
 eg 1,  $\infty$   
 2, 2

## Comparing prob dists

total variation distance

$T(p, q) = \frac{1}{2} \|p - q\|_1 = \max_{\|x\|_\infty \leq 1} \frac{1}{2} \langle p - q, x \rangle = \max_{S \subseteq \Omega} |p(S) - q(S)|$

## fidelity

$\langle \sqrt{p}, \sqrt{q} \rangle = \sum_x \sqrt{p(x)q(x)} =: F(p, q)$

$1 - F \leq T \leq \sqrt{2(1 - F)}$

$F(p_1 \otimes p_2, q_1 \otimes q_2) = F(p_1, q_1) F(p_2, q_2)$

$\Rightarrow 1 - T(p^{\otimes n}, q^{\otimes n}) \sim e^{-cn}$

## 7. states

pure states

$$\| |\alpha\rangle - |\beta\rangle \|_2 = \sqrt{2(1 - \operatorname{Re} \langle \alpha | \beta \rangle)}$$

not great.

$$F(\alpha, \beta) = |\langle \alpha | \beta \rangle| \quad \text{avoids phase dependence}$$

mixed states

$$T(\rho, \sigma) = \frac{1}{2} \|\rho - \sigma\|_1 = \max_{0 \leq M \leq I} \operatorname{tr} M(\rho - \sigma)$$

properties

$$T(V\rho V^\dagger, V\sigma V^\dagger) = T(\rho, \sigma)$$

$$T(\mathcal{E}(\rho), \mathcal{E}(\sigma)) \leq T(\rho, \sigma) \quad \text{pset}$$

note that saturates

$$\mathcal{E}(\rho) = \operatorname{tr}(M\rho) |0\rangle\langle 0| + \operatorname{tr}((\mathbb{I} - M)\rho) |1\rangle\langle 1|$$

this

fidelity of mixed states

$$F(\rho, \sigma) = \|\sqrt{\rho} \sqrt{\sigma}\|_1 = \operatorname{tr} \sqrt{\sqrt{\rho} \sigma \sqrt{\rho}}$$

pset

$$1 - F \leq T \leq \sqrt{1 - F^2}$$

$$\text{and } F(\mathcal{E}(\rho), \mathcal{E}(\sigma)) \geq F(\rho, \sigma)$$

$$0 \leq F \leq 1$$

$\arccos F$  is a metric

Uhlmann's theorem

$$F(\rho, \sigma) = \max |\langle \alpha | \beta \rangle| \text{ over } |\alpha\rangle_{AB}, |\beta\rangle_{AB} \text{ s.t. } \rho_A = \rho, \sigma_B = \sigma$$

by purification uniqueness, we can assume  $|\alpha\rangle$  is fixed

$$|\phi^\rho\rangle = (\sqrt{\rho} \otimes \mathbb{I}) |\Gamma\rangle \quad |\Gamma\rangle = \sum_i |i\rangle \otimes |i\rangle$$

PF

$$|\alpha\rangle = |\phi^\rho\rangle \quad |\beta\rangle = (\mathbb{I} \otimes U) |\phi^\rho\rangle$$

$$\text{RHS} = \max_U |\langle \phi^\rho | \mathbb{I} \otimes U | \phi^\rho \rangle|$$

$$= \max_U |\langle \Gamma | \underbrace{(\sqrt{\rho} \otimes \mathbb{I})(\mathbb{I} \otimes U)(\sqrt{\rho} \otimes \mathbb{I})}_{\sqrt{\rho} \sigma \otimes U} | \Gamma \rangle|$$

$$= \max_U |\operatorname{tr} \sqrt{\rho} \sigma U^\dagger}| \stackrel{*}{=} \|\sqrt{\rho} \sqrt{\sigma}\|_1 = F$$

Justifying  $\star$   $S_1/S_\infty$  duality

$$A = VDW^T \quad \max_U |\text{tr} AU| = |\text{tr} DW^TUV|$$
$$= \sum_i D_{ii} \underbrace{(W^TUV)_{ii}}_{\leq 1} \leq \|A\|_1$$

achieved if  $W^TUV = I$

No-go for noisy bit commitment

After commit states are  $|\psi_0\rangle_{AB}$  or  $|\psi_1\rangle_{AB}$

$$T(\psi_0^B, \psi_1^B) \leq \epsilon$$

$$\Rightarrow F(\psi_0^B, \psi_1^B) \geq 1 - \epsilon$$

$$\Rightarrow \exists U \text{ st. } F(\underbrace{(U \otimes I)\psi_0}_{|\psi_{\text{fake}}\rangle}, \psi_1) \geq 1 - \epsilon$$

$$T(\psi_{\text{fake}}, \psi_1) \leq \sqrt{2\epsilon}$$

$$T(\text{reveal}(\psi_{\text{fake}}), \text{reveal}(\psi_1)) \leq \sqrt{2\epsilon}$$

What do  $T$  &  $F$  measure?

$$\begin{pmatrix} 1 \\ 0 \end{pmatrix} \text{ vs } \begin{pmatrix} 1-\epsilon \\ \epsilon \end{pmatrix} \quad \text{or} \quad \begin{pmatrix} 1/2 \\ 1/2 \end{pmatrix} \text{ vs } \begin{pmatrix} 1/2+\epsilon \\ 1/2-\epsilon \end{pmatrix}$$

$$\text{or } |0\rangle \text{ vs } \cos\theta|0\rangle + \sin\theta|1\rangle$$